

# Spatial networks with random connections

Carl P. Dettmann (University of Bristol)

with Justin P. Coon (Oxford) and Orestis Georgiou (Toshiba)



# Outline

- Introduction to spatial networks
- Random geometric graphs
- Random connection models: Phys. Rev. E **93**, 032313 (2016)
- K-connectivity: EPL **103**, 28006 (2013)
- Anisotropy: Trans. Wireless Commun. **13**, 4534 (2014)

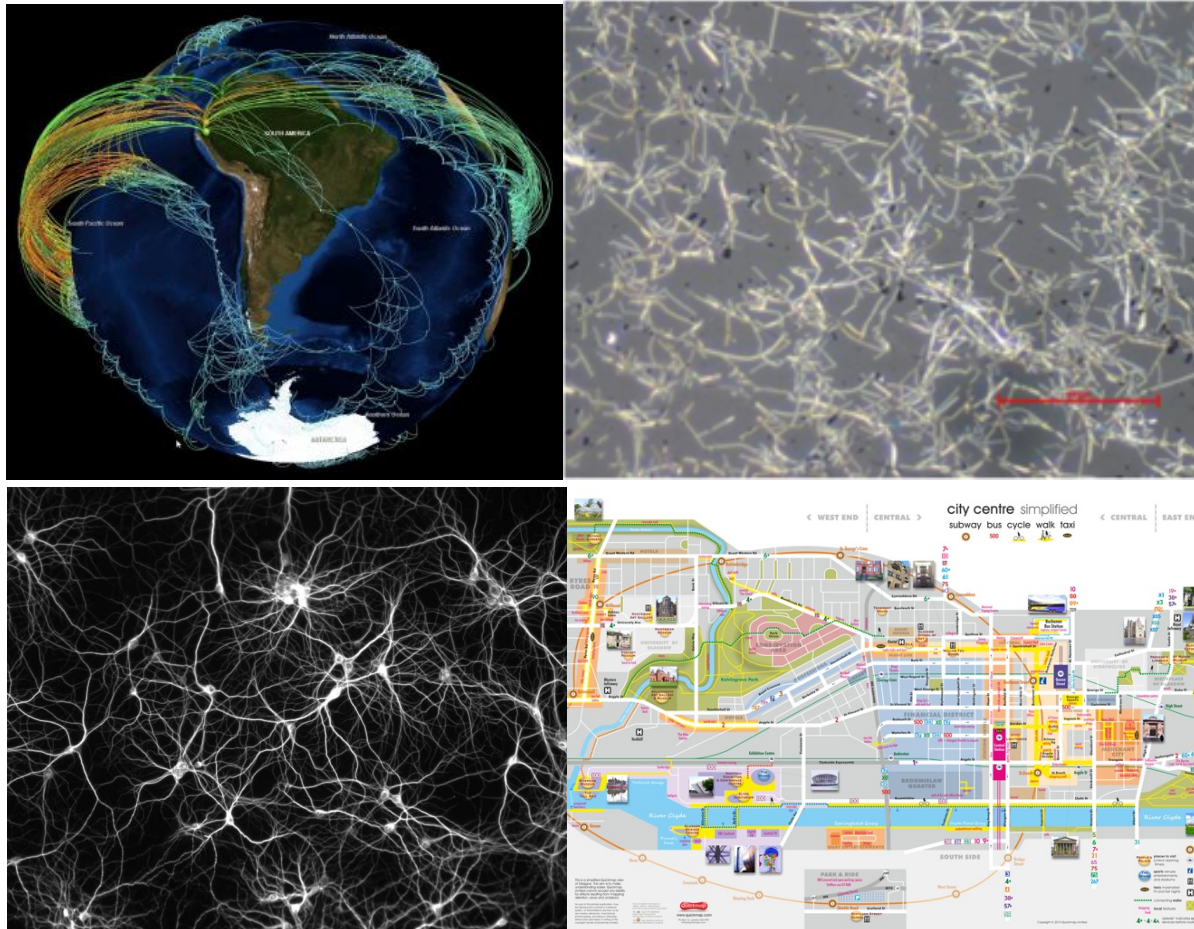
# Networks

A graph  $G = (V, E)$  is a set of vertices  $V$  (“nodes”) together with a subset of unordered pairs (“edges” or “links,”  $E$ ).

- Node degree  $k_i$ : Number of nodes linked to node  $i$ .
- Mean degree  $\mathcal{K} = \frac{\sum_i k_i}{\sum_i 1}$ .
- Path: Sequence of nodes such that adjacent nodes are linked by an edge.
- (Fully) connected: A graph for which there is a path between any pair of nodes. Implies that all  $k_i > 0$ .
- (Fully) connection probability  $P_{fc}$ : Probability of (full) connection in a random graph model.

# Spatial networks

Nodes (and sometimes links) have a location in space.



# Random geometric graphs

Introduced in 1961 by E. N. Gilbert:

*Recently random graphs have been studied as models of communications networks. Points (vertices) of a graph represent stations; lines of a graph represent two-way channels. . . . To construct a random plane network, first pick points from the infinite plane by a Poisson process with density  $D$  points per unit area. Next join each pair of points by a line if the pair is separated by distance less than  $R$ .*

Then:

**Communications networks** Many authors, since 1980s

**Connectivity threshold** Penrose (1997), Gupta & Kumar (1999)

**Books/reviews:**

**Meester & Roy (1996)** Continuum percolation

**Penrose (2003)** Random geometric graphs

**Franceschetti & Meester (2008)** Random networks for communication

**Walters (2011)** Random geometric graphs

**Barthélemy (2011)** Spatial networks

**Haenggi (2012)** Stochastic geometry for wireless networks

## Wireless network considerations

**Mesh architectures** Multihop connections rather than direct to a base station: Reduces power requirements, interference, single points of failure.

**Random node locations** In many applications (sensor, vehicular, swarm robotics, disaster recovery, ...) device locations are unplanned and/or mobile.

**Network characteristics** Full connectivity, k-connectivity (resilience), algebraic connectivity (synchronisation), betweenness centrality (importance, overload).

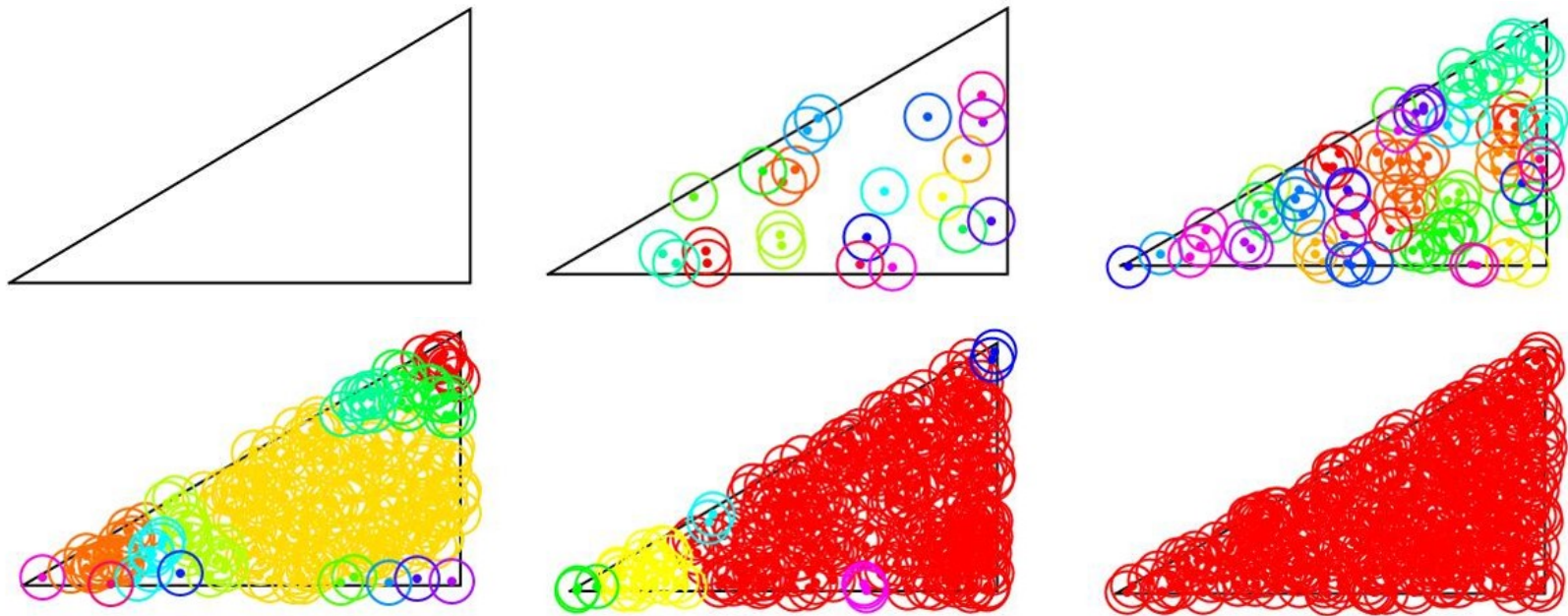
Useful extensions:

**Random connection models** Extra randomness: Link with (iid) probability  $H(r) \in [0, 1]$ , a function of mutual distance  $r$ .

**Anisotropy** Orientations as well as positions.

**Line of sight condition** Impenetrable and/or reflecting boundaries: Particular relevance to millimetre waves.

## Example: A triangle



Isolated nodes occur mostly near the corners...

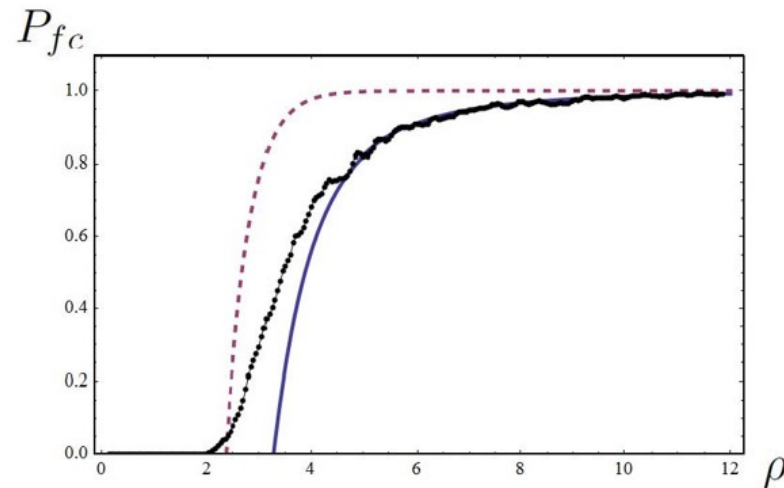
## Dependence on density and geometry

We see two main transitions as density increases:

**Percolation** Formation of a cluster comparable to system size:  
Largely independent of geometry.  $\mathcal{K} = 4.5122\dots$

**Connectivity** All nodes connected in multi-hop fashion:  
Strongly dependent on geometry.  $\mathcal{K} \approx \ln N$ .

What is the full connection probability as a function of density and geometry?





## Previous results

Mathematically rigorous results are for  $N \rightarrow \infty$ , with an appropriate scaling of at least two of  $r_0$ ,  $\rho$  and the system size  $L$ .

For the random geometric graph in dimension  $d \geq 2$ , it was shown by Penrose, and by Gupta & Kumar, that the  $r_0$  threshold for **connectivity** is almost always the same as for **isolated nodes**.

In turn, isolated nodes are local events, so described by a limiting Poisson process: The probability of a node having degree  $k$  is given by

$$P(k) = \frac{\mathcal{K}^k}{k!} e^{-\mathcal{K}}$$

where  $\mathcal{K}$  is the mean degree, equal to  $\rho\pi r_0^2$  for the 2D RGG. This leads to

$$P_{fc} \approx \exp \left[ -\rho V e^{-\rho\pi r_0^2} \right]$$

where  $V$  is the “volume” (ie area) of the domain.

At fixed probability and connection range,  $V$  increases exponentially with  $\rho$

## Random connection model

Penrose (2015) gives proofs for many  $H(r)$  of compact support as  $N \rightarrow \infty$ ; we assume true more generally

$$P_{fc} \approx \exp \left[ - \int \rho e^{-\rho \int H(r_{12}) d\mathbf{r}_2} d\mathbf{r}_1 \right]$$

where  $\rho$  is the density,  $H(r)$  is the iid probability of connection between nodes with mutual distance  $r$  and the integrals are over the domain  $\mathcal{V} \subset \mathbb{R}^d$ .

We want to approximate  $P_{fc}$  for finite  $\rho$ , taking into account boundaries.

**Open problem:** 1D, eg vehicular networks!

## Specific random connection models

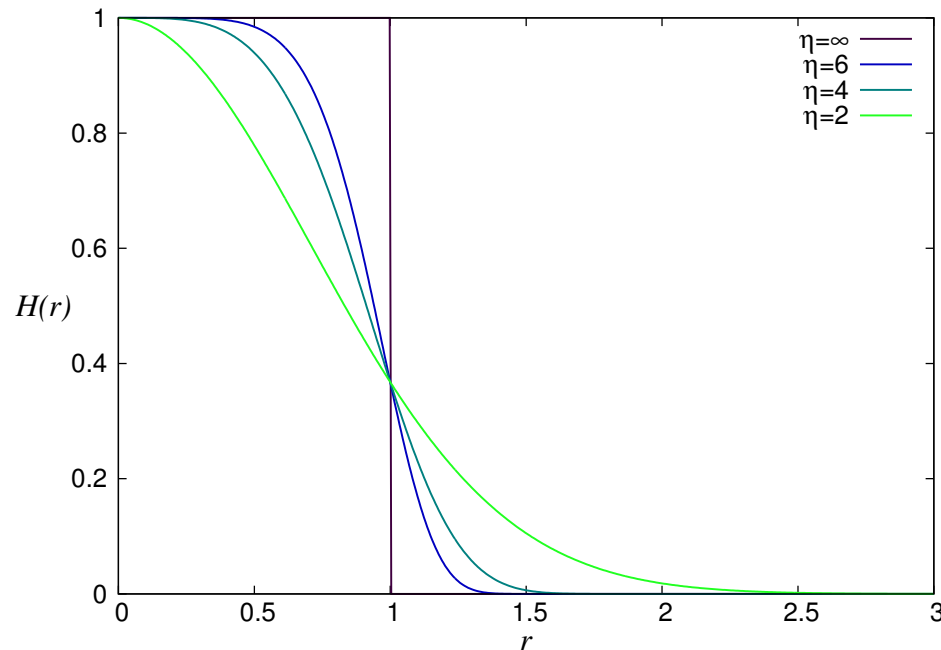
The connection function is the complement of the outage probability,

$$H(r) = \mathbb{P}(\log_2(1 + SNR |h|^2) > R_0)$$

neglecting interference, with  $SNR \propto r^{-\eta}$ , path loss exponent  $\eta \in [2, 6]$ , rate  $R_0$ . Simplest is Rayleigh fading (diffuse signal), for which the channel gain  $|h|^2$  is exponentially distributed, giving

$$H(r) = \exp[-(r/r_0)^\eta]$$

Similar, though more involved: MIMO, Rician (specular plus diffuse), ...



## Connectivity and boundaries

For large  $\rho$ , dominated by the regions of small connectivity mass

$$M(\mathbf{r}_2) = \int H(r_{12}) d\mathbf{r}_1$$

Exactly on the boundary, this is given by

$$M_B = H_{d-1} \omega_B$$

where

$$H_m = \int_0^\infty H(r) r^m dr$$

is the  $m$ th moment, and  $\omega_B$  is the (solid) angle associated with the boundary component  $B$ , eg  $\pi/2$  for a right angled corner,  $\pi$  for an edge.

Analysing the vicinity of boundaries more carefully...

## General formula

$$P_{fc} = \exp \left[ - \sum_B \rho^{1-i_B} G_B V_B e^{-\rho \omega_B H_{d-1}} \right]$$

where  $i_B$  is the boundary codimension,  $V_B$  is its  $d - i$  dimensional volume, and  $G_B$  is the geometrical factor

$G_B$	$i = 0$	$i = 1$	$i = 2$	$i = 3$
$d = 2$	1	$\frac{1}{2H_0}$	$\frac{1}{H_0^2 \sin \omega}$	
$d = 3$	1	$\frac{1}{2\pi H_1}$	$\frac{1}{\pi^2 H_1^2 \sin(\omega/2)}$	$\frac{4}{\pi^2 H_1^3 \omega \sin \omega}$

where the 3D corner has a right angle.

**Curved boundaries?** To leading order, modification of the exponential but not the geometrical factor:

$$P_{2,1} = \dots e^{-\rho(\pi H_1 - \kappa H_2)}$$

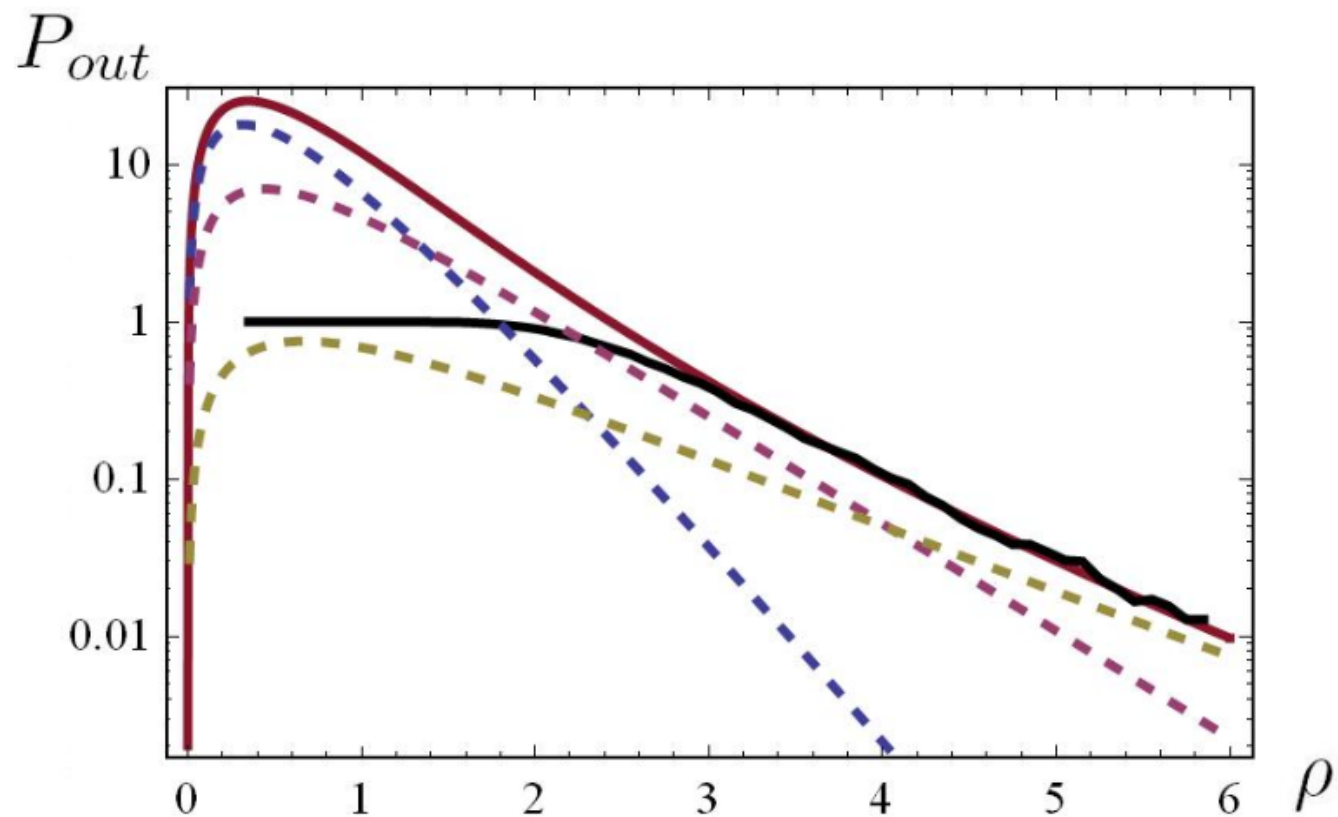
$$P_{3,1} = \dots e^{-\pi \rho(2H_2 - \kappa H_3)}$$

where  $\kappa$  is (mean) curvature.

## Example: A square

The previous formula gives

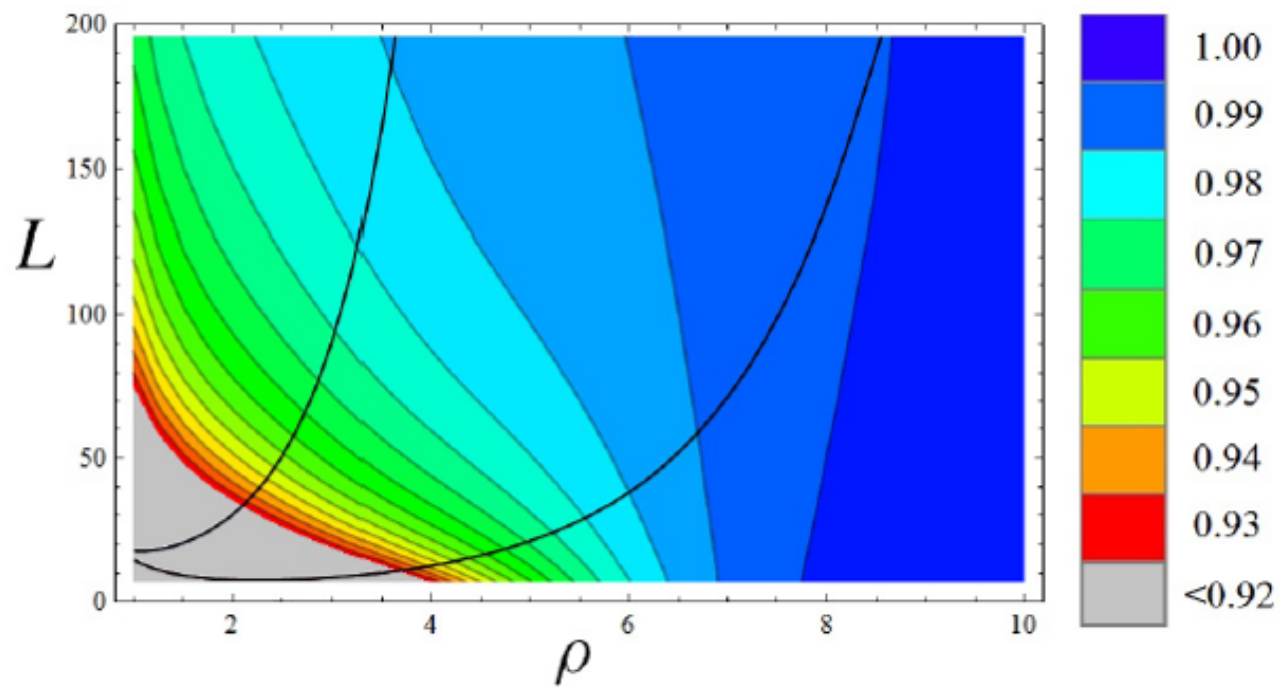
$$1 - P_{fc} \approx L^2 \rho e^{-\pi \rho} + \frac{4L}{\sqrt{\pi}} e^{-\frac{\pi \rho}{2}} + \frac{16}{\pi \rho} e^{-\frac{\pi \rho}{4}}$$



## Phase diagram

Testing convergence of

$$\frac{1 - P_{fc}}{\sum_B \dots}$$

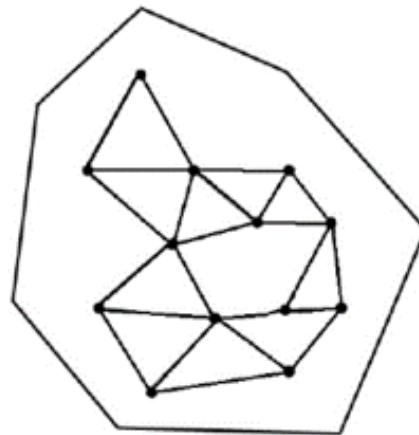


## K-connectivity

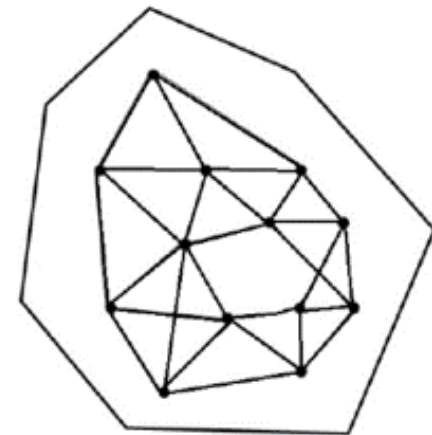
A network is (vertex)  $k$ -connected if any  $k - 1$  nodes can be removed and it remains connected. It is a useful measure of **network resilience**.



1-connected



2-connected



3-connected

Vertex connectivity  $\leq$  Edge connectivity  $\leq$  Minimum degree



## Minimum degree

Assume independence ...

- For each node, degree is Poisson:

$$P_i(k) \approx \frac{\mathcal{K}_i^k}{k!} e^{-\mathcal{K}_i}$$

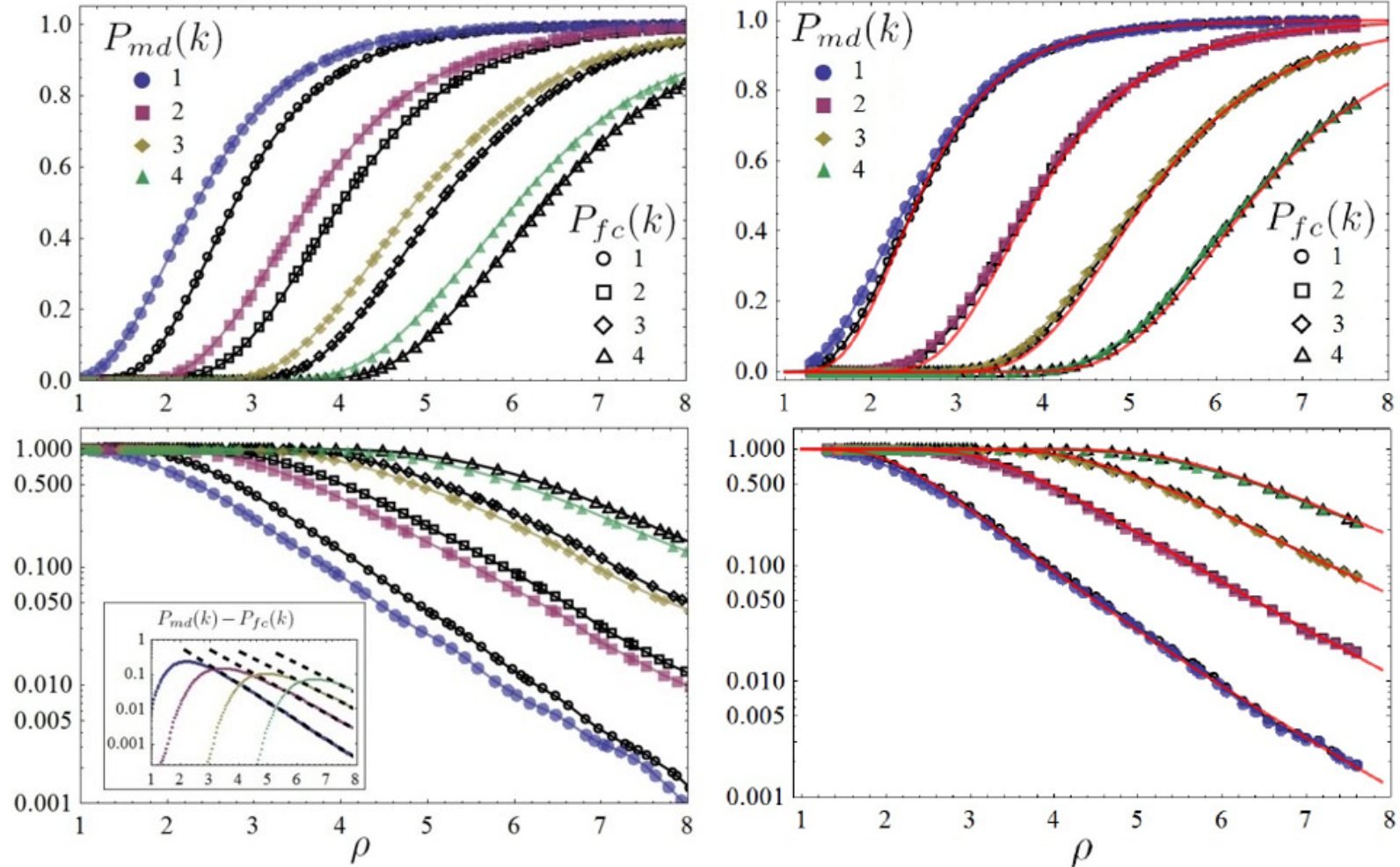
- Node degrees are independent:

$$P_{md}(k) \approx \left[ 1 - \sum_{m=0}^{k-1} \frac{\rho^m}{m!} \frac{1}{V} \int_{\mathcal{V}} M_H^m(\mathbf{r}_i) e^{-\rho M_H(\mathbf{r}_i)} d\mathbf{r}_i \right]^N$$

# Numerical results

Hard

Soft



Random connections: Minimum degree is a better proxy for  $k$ -connectivity.

Why? Connections are less correlated in the random model.

## Anisotropic connections

- Angle-dependent transmit and receive gains:

$$H(r, \phi, \theta_T, \theta_R) = \exp \left( -\frac{\beta r^\eta}{G_T(\phi - \theta_T)G_R(\phi + \pi - \theta_R)} \right)$$

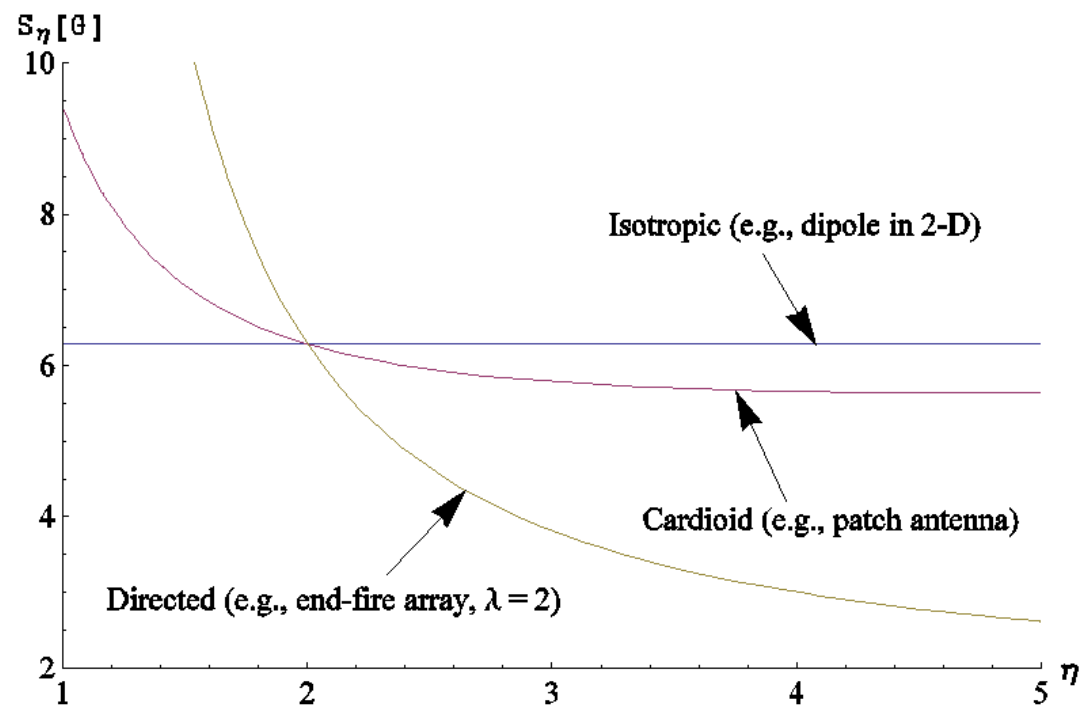
- Fix total power per node

$$\int_0^{2\pi} G_T(\phi) d\phi = \int_0^{2\pi} G_R(\phi) d\phi = 2\pi$$

- Connectivity mass is now

$$M = \frac{1}{2\pi} \int r H(r, \phi, \theta_T, \theta_R) dr d\phi d\theta_R$$

## Transition at $\eta = d$

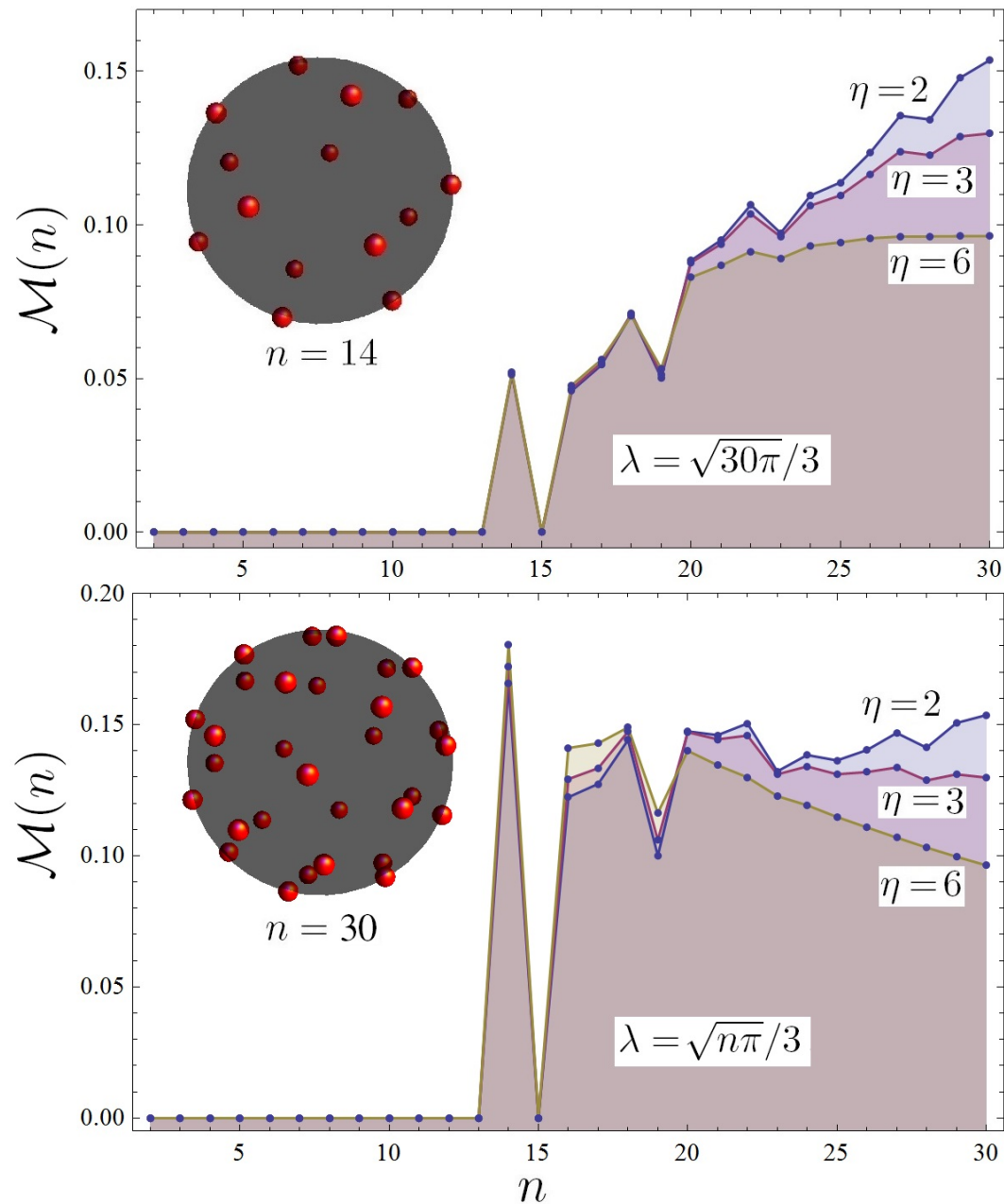


## Anisotropy and boundaries

- Homogeneous case:
  - Path loss exponent  $\eta > d$ : Isotropic optimal
  - Path loss exponent  $\eta < d$ : Delta spike(s) optimal
- With boundaries, for  $\eta < d$ , trade-off between system size/shape and number/width of spikes. Examples:
  - Square, best to have a multiple of 4 spikes.
  - Cube ...

# Cube optimal pattern

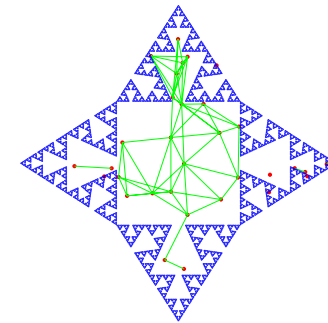
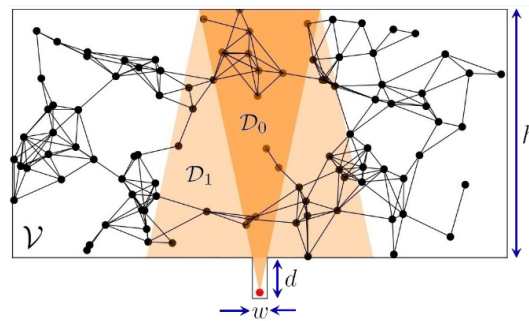
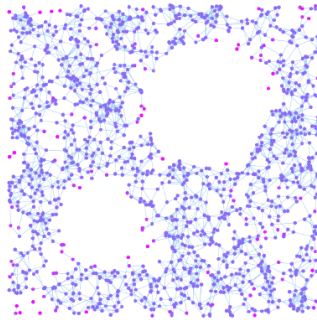
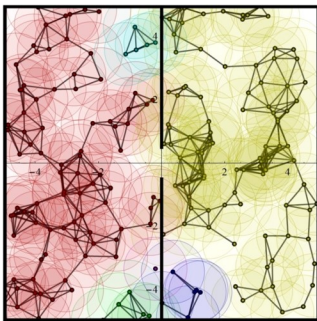
14 spikes: Gyroelongated hexagonal bipyramid!



# Outlook

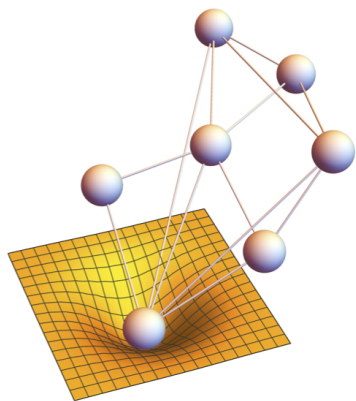
Random connection models are more realistic and have smoother properties.

Other results/in progress: Non-convex domains, betweenness, interference, nonuniform, mobility, spectrum . . .



Connection functions for other spatial networks?

Very long range connections?



SPATIALLY  
EMBEDDED  
NETWORKS

EPSRC

Engineering and Physical Sciences  
Research Council